

A Mathematical Challenge: Bookings and Technical Capacities in the European Entry-Exit Gas Market

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Discrete Optimization, FAU Erlangen-Nürnberg

Energy Seminar Edinburgh, May 7 2020





Perfect Competition

Multilevel Entry-Exit Gas Market Model

A four level model for the European Entry-Exit Gas Market¹:



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A four level model for the European Entry-Exit Gas Market¹:

Maximize social welfare



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Level 4: Cost-optimal transport through the network by the TSO



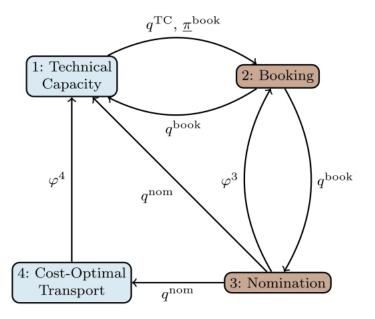


Figure: A Multi-level Gas Market Model¹



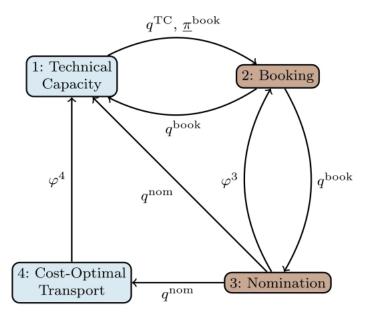


Figure: A Multi-level Gas Market Model¹

Main Goal: Decouple trading and transport using technical capacities



Structure of Bilevel Model

Upper Level (TSO):

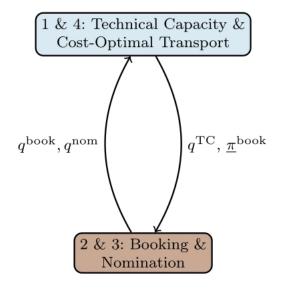


Figure: Bilevel Reformulation¹



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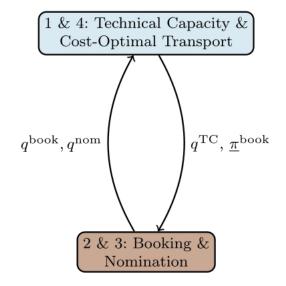


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Upper Level (TSO):

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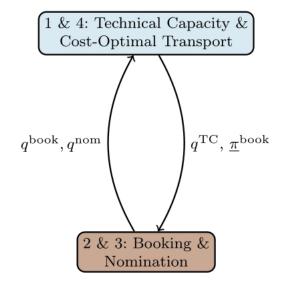


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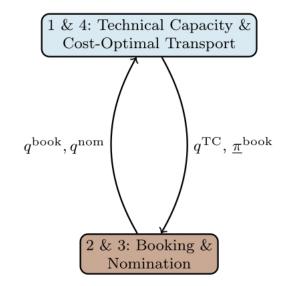


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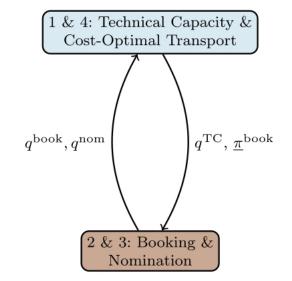


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Main Challenge: technical capacities



Nonlinear Adjustable Robustness

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$$\begin{aligned} \pi_u - \pi_v &= \psi_a \left(d_a, q_a \right), & a = (u, v) \in A, \\ \sum_{a \in \delta^{\text{out}}(v)} q_a - \sum_{a \in \delta^{\text{in}}(v)} q_a = q_v^{\text{nom}}, & v \in V, \\ \pi_v &\leq \pi_v \leq \bar{\pi}_v, & v \in V, \end{aligned}$$

is feasible.



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is feasible.

No standard techniques of robust optimization are applicable



Joint Work Robinius, Schewe, Schmidt, Stolten, and Welder²



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Potential Loss Function

Here, $\psi_a(d_a, \cdot) : \mathbb{R} \to \mathbb{R}$ is a *potential function* that is continuous, strictly increasing, and odd w.r.t the second argument.

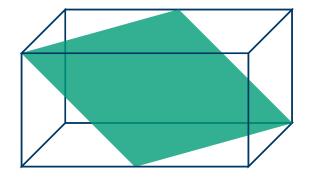


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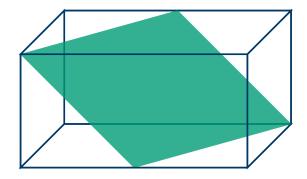


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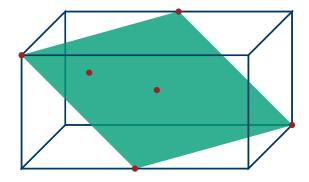


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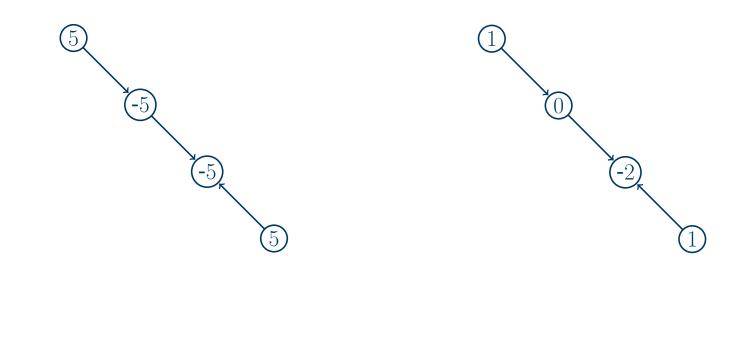




Greedy Approach

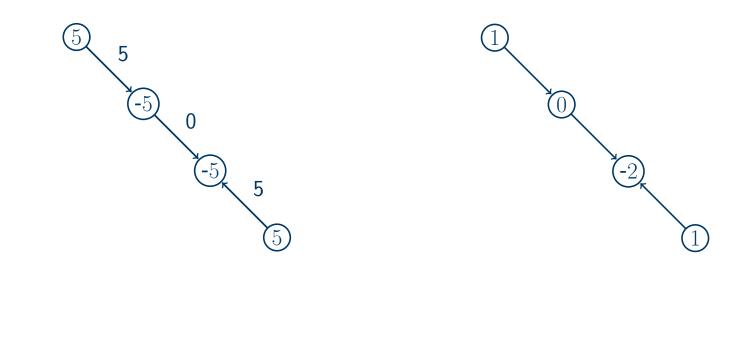


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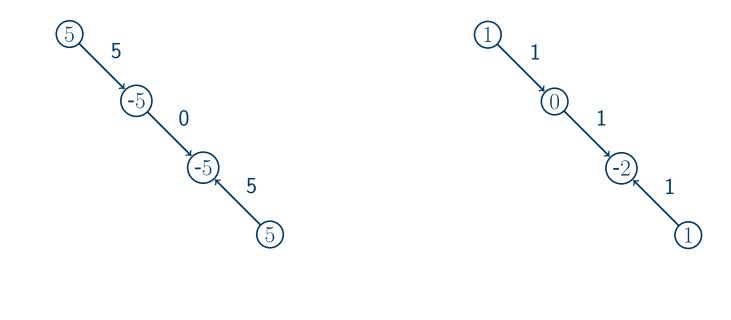


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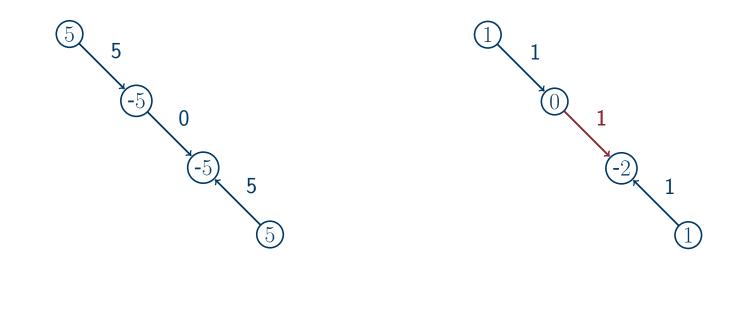


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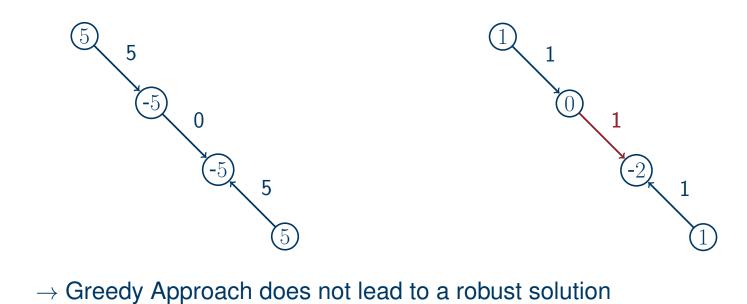


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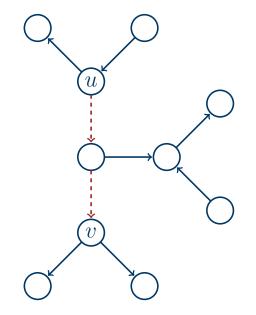
Technical capacities are feasible if and only if these special nominations are feasible.



Structure of Special Nominations

Special Nominations

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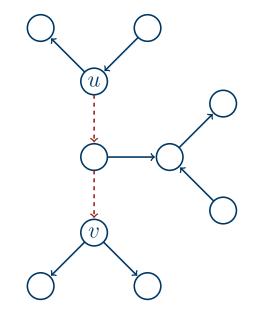


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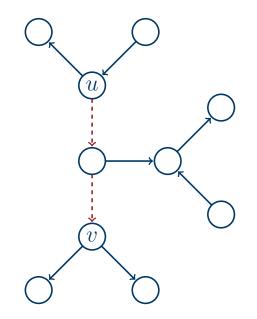
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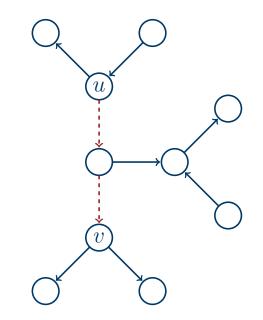
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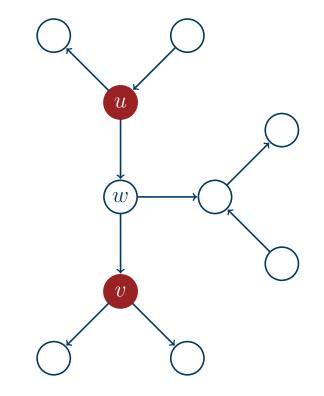
Lemma (Existence of Special Nominations)

For each entry u and exit v a special nomination exists: $B_{u,v} \neq \emptyset$ and it can be computed in polynomial time by an LP.



Lemma

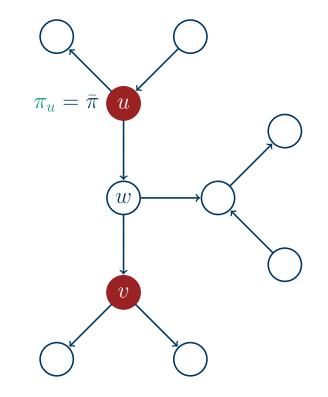
- the pressure drop from u to v in $q_{u,v}^{\text{nom}}$ is larger than in q^{nom} and
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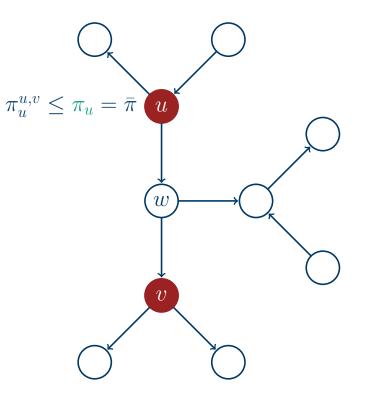
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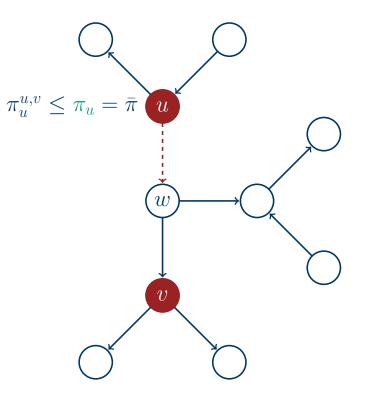
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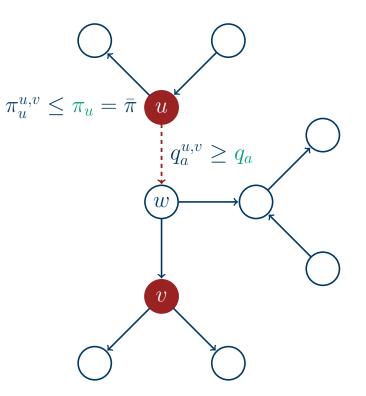
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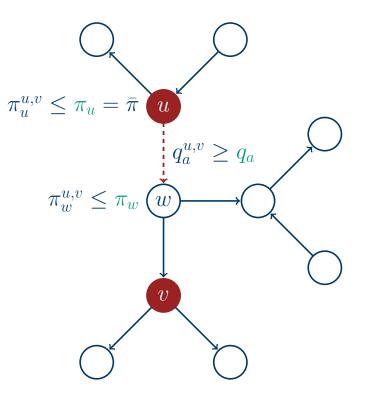
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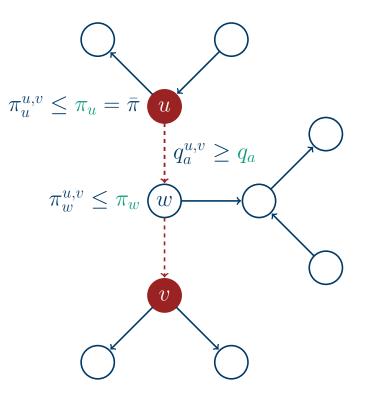
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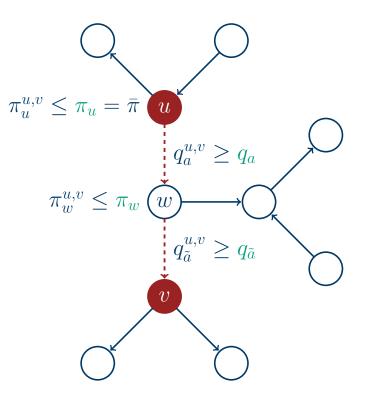
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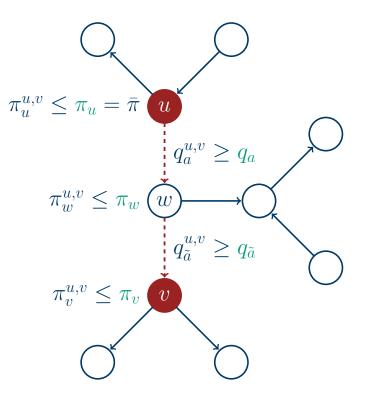
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Diameter Selection of Hydrogen Networks

Detailed Setting

- Tree G = (V, A) with passive arcs
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- Technical capacities q^{TC} , here demand prediction

$$B := \left\{ q^{\text{nom}} \colon \sum_{v \in V} \sigma_u q_v^{\text{nom}} = 0, \ 0 \le q_v^{\text{nom}} \le q^{\text{TC}}, \ v \in V \right\}$$



Modeling Nonlinear Model

Diameter Sizing Model

$$\begin{split} \min_{x,q,\pi} & \sum_{a \in A} \sum_{d \in D_a} c_{a,d} x_{a,d} \\ \text{s.t.} & \pi_{q_u^{\text{nom}}} - \pi_{q_v^{\text{nom}}} = \psi_a \left(\sum_{d \in D_a} d x_{a,d}, q_{q_a^{\text{nom}}} \right), \quad a = (u,v) \in A, \ q^{\text{nom}} \in B, \\ & \sum_{a \in \delta^{\text{out}}(v)} q_{q_a^{\text{nom}}} - \sum_{a \in \delta^{\text{in}}(v)} q_{q_a^{\text{nom}}} = q_v^{\text{nom}}, \qquad v \in V, \ q^{\text{nom}} \in B, \\ & \pi_v \leq \pi_{q_v^{\text{nom}}} \leq \bar{\pi}_v, \qquad v \in V, \ q^{\text{nom}} \in B, \\ & \sum_{d \in D_a} x_{a,d} = 1, \qquad a \in A, \\ & x_{a,d} \in \{0,1\}, \qquad a \in A, \ d \in D_a. \end{split}$$



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Details of Instance

Instance

- Realistic hydrogen network of Eastern Germany planned by Forschungszentrum Jülich (IEK-3)
- 1 Entry, 1 storage, and 745 exits
- 28 diameters in range 0.1063 – 1.536 m
- Upper pressure bound $\bar{\pi} =$ 95 bar

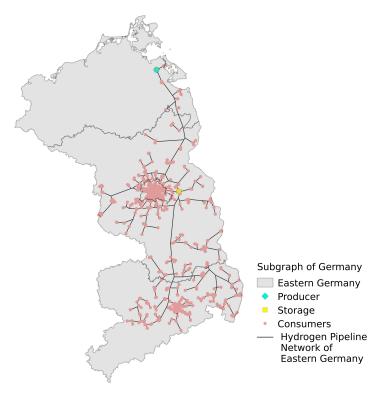


Figure: Hydrogen Network²

²Robinius, Schewe, Schmidt, Stolten, Thürauf, and Welder, "Robust Optimal Discrete Arc Sizing for Tree-Shaped Potential Networks".



Results of Instance: Refueling Station without Local Onsite Storage

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- Reduced nomination set *B** contains 25 nominations
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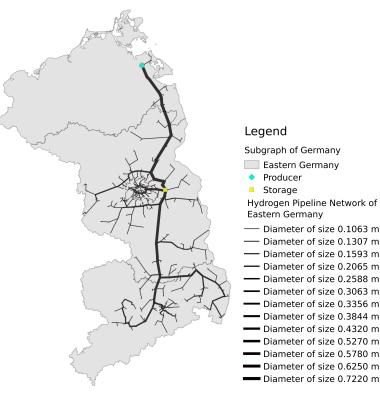
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Vector q^{TC} are feasible technical capacities if and only if

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$$\varphi_{w_1w_2}(q^{\mathrm{TC}}) \le \pi_{w_1}^+ - \pi_{w_2}^-$$
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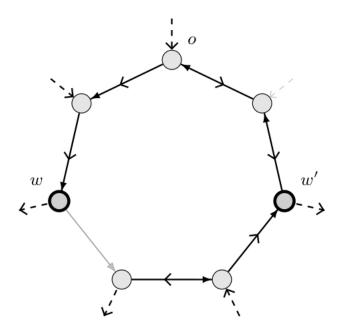


Figure: Different flow-meeting points w and w'.⁴

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Cycles

Results of trees cannot be used for cycles

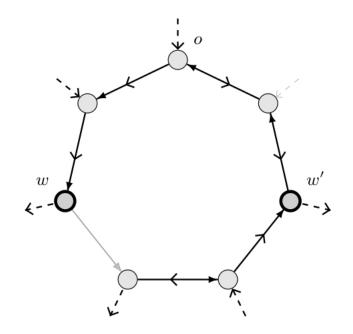


Figure: Different flow-meeting points w and $w^{\prime}.^4$

⁴Labbé, Plein, Schmidt, and Thürauf, *Deciding Feasibility of a Booking in the European* Gas Market on a Cycle is in P.



Joint work with M. Labbé, F. Plein, M. Schmidt⁴

Cycles

Results of trees cannot be used for cycles

Why are cycles so bad?

No cyclic flows are possible

Different flow-meeting points

The potential drop in a cycle sums up to zero

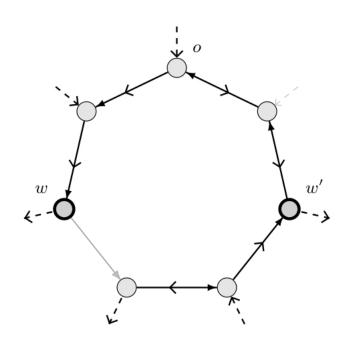


Figure: Different flow-meeting points w and w'.⁴

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Theorem²

Vector q^{TC} are feasible technical capacities if and only if

$$\varphi_{w_1w_2}(q^{\mathrm{TC}}) \leq \pi_{w_1}^+ - \pi_{w_2}^- \text{ for all } (w_1, w_2) \in V^2$$

$$\varphi_{w_1w_2}(q^{\mathrm{TC}}) \coloneqq \max_{q^{\mathrm{nom}}, q} \quad \pi_{w_1} - \pi_{w_2}$$

s.t.
$$\sum_{a \in \delta^{\mathrm{out}}(v)} q_a - \sum_{a \in \delta^{\mathrm{in}}(v)} q_a = \sigma_u q_u^{\mathrm{nom}}, \qquad u \in V,$$

$$\pi_u - \pi_v = \Lambda_a q_a |q_a|, \qquad a = (u, v) \in A,$$

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Theorem

It always exists an optimal solution with

- a single flow-meeting point
- a *certain* combinatorial structure and only polynomial many different possibilities for this structure exist.



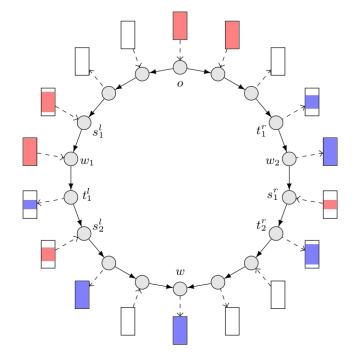


Figure: Properties of an Optimal Solution that Maximizes the Potential Difference between w_1 and w_2 .⁴

⁴Labbé, Plein, Schmidt, and Thürauf, *Deciding Feasibility of a Booking in the European* Gas Market on a Cycle is in P.



Sketch Algorithm

Result: Maximum potential difference between w_1 and w_2 . foreach Combinatorial Structure (polynomially many) do



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We can decide the feasibility of technical capacities in $O((\log(V) + \tau) + |V^+|^5|V^-|^4)$.



Complexity of Technical Capacities

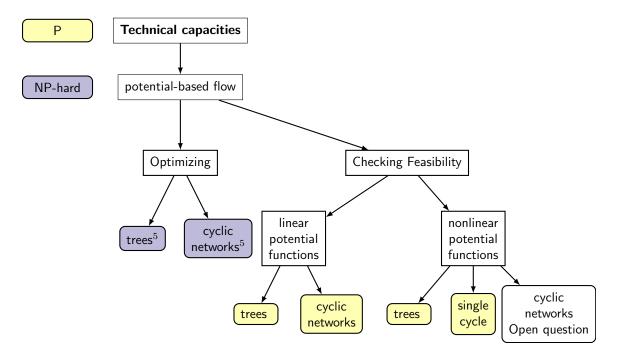


Figure: Overview of known complexity results for technical capacities and potential-based flows.

⁵Schewe, Schmidt, and Thürauf, *Computing Technical Capacities in the European Entry-Exit Gas Market is NP-Hard*.



Future Research

Solve the nonlinear bilevel model for the Entry-Exit gas market system with nonlinear potential-based flows for trees.



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Finish my PhD



Journal Articles

Martin Robinius, Lars Schewe, Martin Schmidt, Detlef Stolten, Johannes Thürauf, and Lara Welder *Robust Optimal Discrete Arc Sizing for Tree-Shaped Potential Networks* Published in Computational Optimization and Applications (2019): https://link.springer.com/article/10.1007/s10589-019-00085-x

Markus Reuß, Lara Welder, Johannes Thürauf, Jochen Linßen, Thomas Grube, Lars Schewe, Martin Schmidt, Detlef Stolten, and Martin Robinius Modeling hydrogen networks for future energy systems: A comparison of linear and nonlinear approaches Published in International Journal of Hydrogen Energy (2019): https://www.sciencedirect.com/science/article/pii/S0360319919338625



Preprints

Martine Labbé, Fränk Plein, Martin Schmidt, and Johannes Thürauf Deciding Feasibility of a Booking in the European Gas Market on a Cycle is in P Preprint: http://www.optimization-online.org/DB HTML/2019/11/7472.html

Lars Schewe, Martin Schmidt, and Johannes Thürauf Computing Technical Capacities in the European Entry-Exit Gas Market is NP-Hard Preprint: http://www.optimization-online.org/DB HTML/2020/01/7576.html

Thank you for your attention